

Functions

Part Two

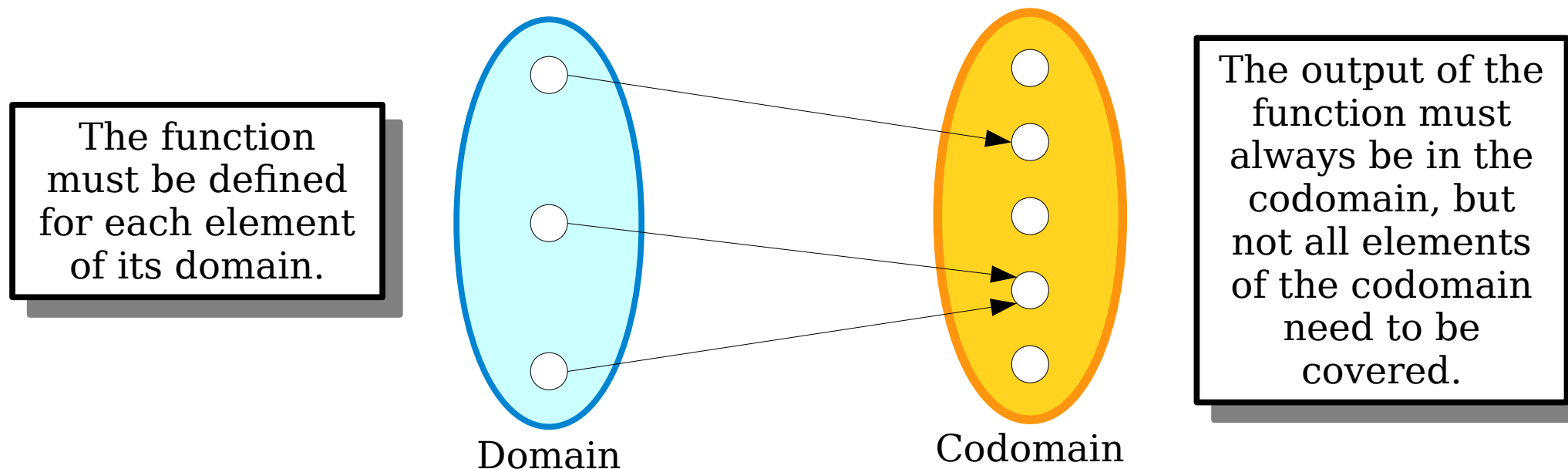
Outline for Today

- ***Recap from Last Time***
 - Where are we, again?
- ***A Proof About Birds***
 - Trust me, it's relevant. \exists
- ***Assuming vs Proving***
 - Two different roles to watch for.
- ***Connecting Function Types***
 - Relating the topics from last time.
- ***Function Composition***
 - Sequencing functions together.

Recap from Last Time

Domains and Codomains

- Every function f has two sets associated with it: its **domain** and its **codomain**.
- A function f can only be applied to elements of its domain. For any x in the domain, $f(x)$ belongs to the codomain.
- We write $f : A \rightarrow B$ to indicate that f is a function whose domain is A and whose codomain is B .



Involutions

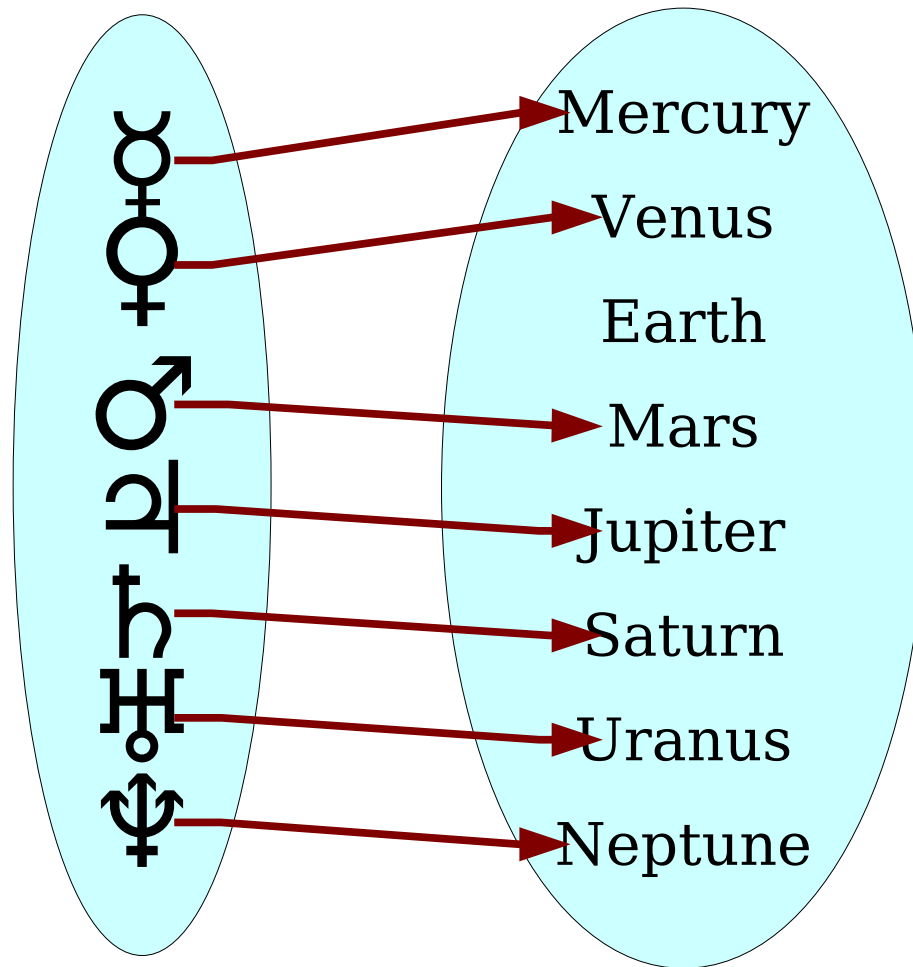
- A function $f : A \rightarrow A$ from a set back to itself is called an ***involution*** if the following first-order logic statement is true about f :

$$\forall x \in A. f(f(x)) = x.$$

(“Applying f twice is equivalent to not applying f at all.”)

Injective Functions

- $\forall a_1 \in A. \forall a_2 \in A. (a_1 \neq a_2 \rightarrow f(a_1) \neq f(a_2))$
- $\forall a_1 \in A. \forall a_2 \in A. (f(a_1) = f(a_2) \rightarrow a_1 = a_2)$



Review: Injective Functions

Injective Functions

Theorem: Let $f : \mathbb{N} \rightarrow \mathbb{N}$ be defined as $f(n) = 2n + 7$.
Then f is injective.

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What does it mean for the function f to be injective?

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$$\forall n_1 \in \mathbb{N}. \forall n_2 \in \mathbb{N}. (f(n_1) = f(n_2) \rightarrow n_1 = n_2)$$

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Write two different sentences that could be the **first** sentence of a **Direct Proof** approach to this proof, one for each of the two definitions of injective—the “assume” step. (remember that *direct proof* is for proving theorems that are implications—in this case that implication is in the definitions of injectivity.)

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Write two different sentences that could be the **second** sentence of a **Direct Proof** approach to this proof, one for each of the two definitions of injective—the “want-to-show” step. (remember that direct proof is for proving theorems that are implications—in this case that implication is in the definitions of injectivity.)

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Therefore, we'll pick arbitrary $n_1, n_2 \in \mathbb{N}$,
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Good exercise: Repeat this proof using the other definition of injectivity!

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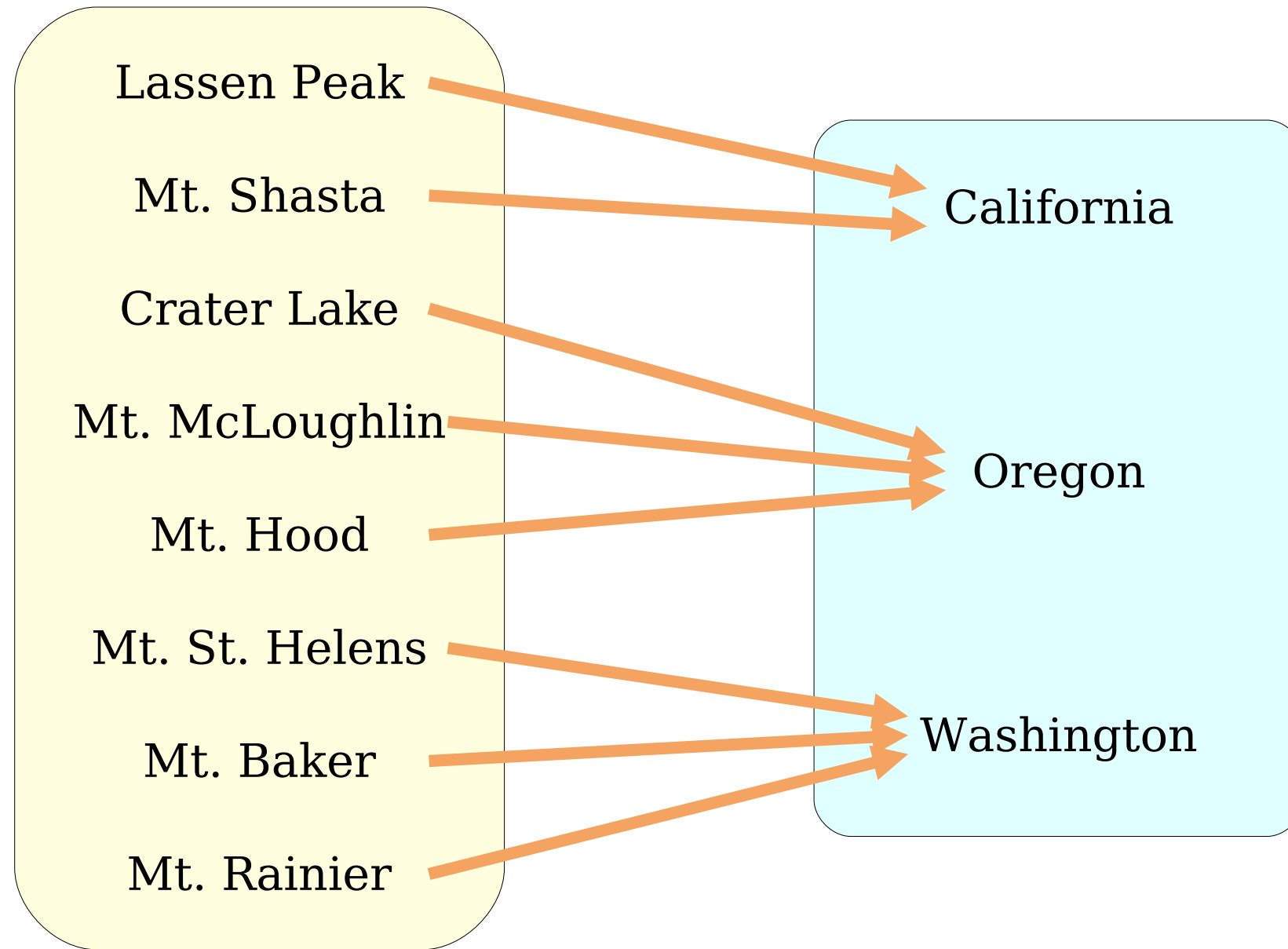
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!! Important style rule !!
This proof contains no
first-order logic syntax
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It's written in plain English,
just as usual.

Another Class of Functions



Surjective Functions

- A function $f : A \rightarrow B$ is called **surjective** (or **onto**) if this first-order logic statement is true about f :

$$\forall b \in B. \exists a \in A. f(a) = b$$

(“For every output, there's an input that produces it.”)

- A function with this property is called a **surjection**.
- How does this compare to our first rule of functions?

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Therefore, we'll choose an arbitrary $y \in \mathbb{R}$, then prove that there is some $x \in \mathbb{R}$ where $f(x) = y$.

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Let $x = y / 2$.

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$\forall x. A$	Have the reader pick an arbitrary x . We then prove A is true for that choice of x .
$\exists x. A$	Find an x where A is true. Then prove that A is true for that specific choice of x .
$A \rightarrow B$	Assume A is true, then prove B is true.
$A \wedge B$	Prove A . Then prove B .
$A \vee B$	Either prove $\neg A \rightarrow B$ or prove $\neg B \rightarrow A$. <i>(Why does this work?)</i>
$A \leftrightarrow B$	Prove $A \rightarrow B$ and $B \rightarrow A$.
$\neg A$	Simplify the negation, then consult this table on the result.

Pop Quiz!
Which row of this
proof techniques
table did we use for
for that proof?

A Proof About Birds



Theorem: If all birds can fly,
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Given the predicates

Bird(*b*), which says *b* is a bird;

Heron(*h*), which says *h* is a heron; and

CanFly(*x*), which says *x* can fly,

translate the theorem into first-order logic.

Go to
[PollEv.com/cs103spr25](https://pollev.com/cs103spr25)

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Theorem: If all birds can fly, then all herons can fly.

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Which makes more sense as the next step in this proof?

1. Consider an arbitrary bird b .
2. Consider an arbitrary heron h .

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Proof: Assume that all birds can fly. We will show that all herons can fly.

Consider an arbitrary heron h . We will show that h can fly. To do so, note that since h is a heron we know h is a bird. Therefore, by our earlier assumption, h can fly. ■

$(\forall b. (Bird(b) \rightarrow CanFly(b))) \rightarrow (\forall h. (Heron(h) \rightarrow CanFly(h)))$

We never introduce a variable b .

We introduce a variable h almost immediately.

Proving vs. Assuming

- In the context of a proof, you will need to assume some statements and prove others.
 - Here, we **assumed** all birds can fly.
 - Here, we **proved** all herons can fly.
- Statements behave differently based on whether you're assuming or proving them.

$(\forall b. (Bird(b) \rightarrow CanFly(b))) \rightarrow (\forall h. (Heron(h) \rightarrow CanFly(h)))$

We never introduce a variable b .

We introduce a variable h almost immediately.

Proving vs. Assuming

- To **prove** the universally-quantified statement

$$\forall x. P(x)$$

we introduce a new variable x representing some arbitrarily-chosen value.

- Then, we prove that $P(x)$ is true for that variable x .
- That's why we introduced a variable h in this proof representing a heron.

$$(\forall b. (Bird(b) \rightarrow CanFly(b))) \rightarrow (\forall h. (Heron(h) \rightarrow CanFly(h)))$$

We never introduce a variable b .

We introduce a variable h almost immediately.

Proving vs. Assuming

- If we **assume** the statement

$$\forall x. P(x)$$

we **do not** introduce a variable x .

- Rather, if we find a relevant value z somewhere else in the proof, we can conclude that $P(z)$ is true.
- That's why we didn't introduce a variable b in our proof, and why we concluded that h , our heron, can fly.

$$(\forall b. (Bird(b) \rightarrow CanFly(b))) \rightarrow (\forall h. (Heron(h) \rightarrow CanFly(h)))$$

We never introduce a variable b .

We introduce a variable h almost immediately.

	To <i>prove</i> that this is true...	
$\forall x. A$	Have the reader pick an arbitrary x . We then prove A is true for that choice of x .	
$\exists x. A$	Find an x where A is true. Then prove that A is true for that specific choice of x .	
$A \rightarrow B$	Assume A is true, then prove B is true.	
$A \wedge B$	Prove A . Then prove B .	
$A \vee B$	Either prove $\neg A \rightarrow B$ or prove $\neg B \rightarrow A$. <i>(Why does this work?)</i>	
$A \leftrightarrow B$	Prove $A \rightarrow B$ and $B \rightarrow A$.	
$\neg A$	Simplify the negation, then consult this table on the result.	

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$\forall x. A$	Have the reader pick an arbitrary x . We then prove A is true for that choice of x .	Initially, <i>do nothing</i> . Once you find a z through other means, you can state it has property A .
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$\exists x. A$	Find an x where A is true. Then prove that A is true for that specific choice of x .	Introduce a variable x into your proof that has property A .
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$A \wedge B$	Prove A . Then prove B .	Assume A . Then assume B .
$A \vee B$	Either prove $\neg A \rightarrow B$ or prove $\neg B \rightarrow A$. <i>(Why does this work?)</i>	Consider two cases. Case 1: A is true. Case 2: B is true.
$A \leftrightarrow B$	Prove $A \rightarrow B$ and $B \rightarrow A$.	Assume $A \rightarrow B$ and $B \rightarrow A$.
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Connecting Function Types

Types of Functions

- We've seen three special types of functions:
 - ***involutions***, functions that undo themselves;
 - ***injections***, functions where different inputs go to different outputs; and
 - ***surjections***, functions that cover their whole codomain.
- ***Question:*** How do these three classes of functions relate to one another?

Theorem: For any function $f : A \rightarrow A$,
if f is an involution, then f is surjective.

$$\underbrace{(\forall x \in A. f(f(x)) = x)}_{\substack{f \text{ is an} \\ \text{involution.}}} \rightarrow \underbrace{(\forall b \in A. \exists a \in A. f(a) = b)}_{\substack{f \text{ is} \\ \text{surjective.}}}$$

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Proof Outline

1. Assume f is an involution.

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$$(\forall x \in A. f(f(x)) = x) \rightarrow (\forall b \in A. \exists a \in A. f(a) = b)$$

Ass
We've said that we need to prove this statement. How do we do that?

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$$(\forall x \in A. f(f(x)) = x) \rightarrow (\forall b \in A. \exists a \in A. f(a) = b)$$

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There's a universal quantifier up front. Since we're proving this, we'll pick an arbitrary $b \in A$.

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Now, we hit an existential quantifier. Since we're proving this, we need to find a choice of $a \in A$ where this is true.

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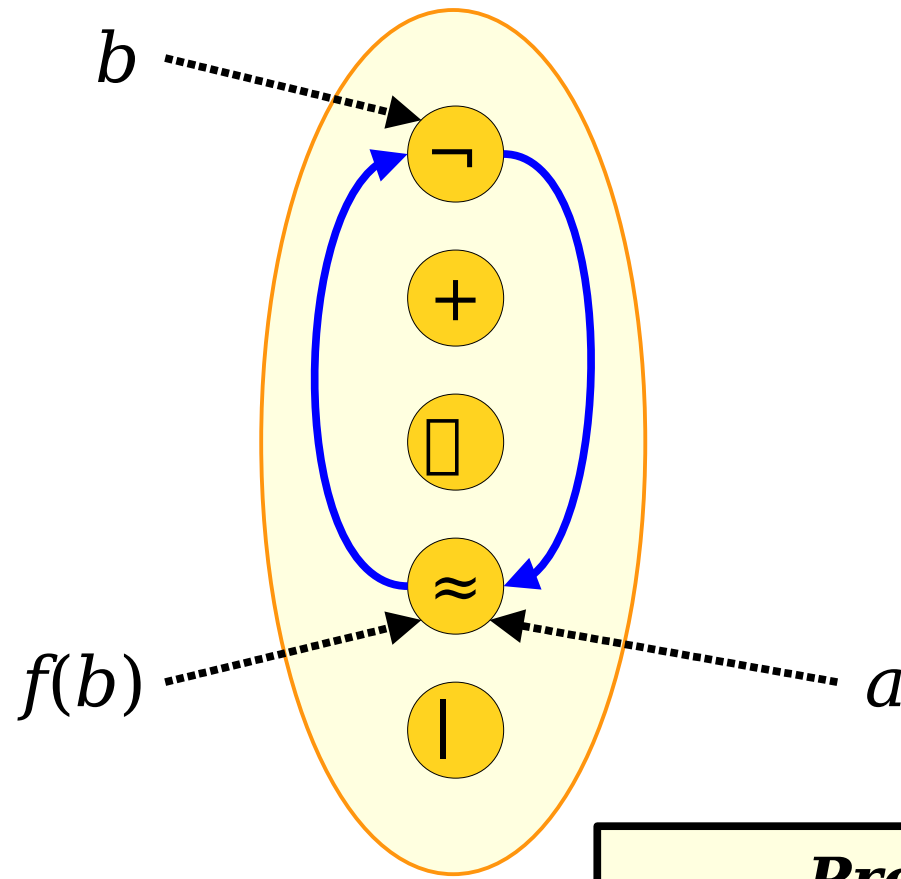
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We need to prove this part.
What does that mean?

Prove
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We now need to prove this implication. But we know how to do that! We assume the antecedent and prove the consequent.

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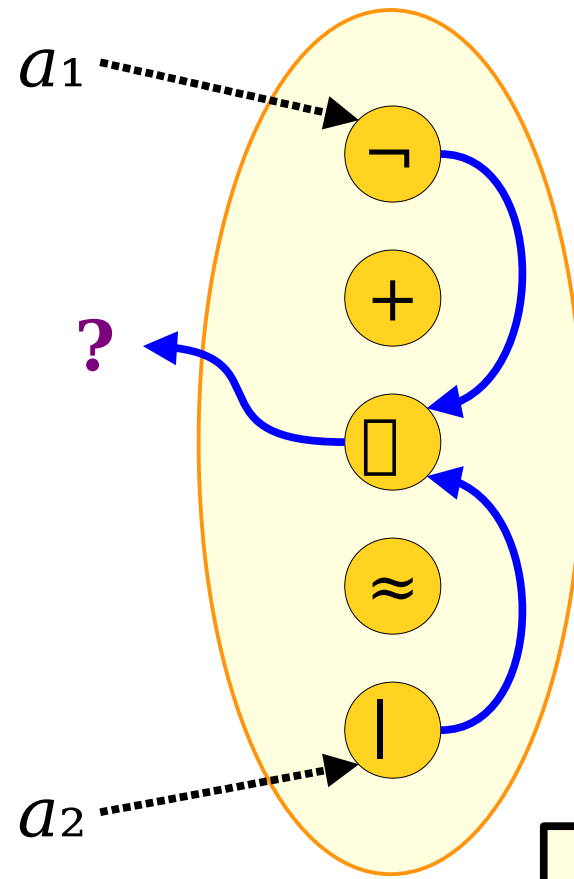
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We've reached a contradiction, so our assumption was wrong.

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Proof: Consider any function $f : A \rightarrow A$ that's an involution. We will prove that f is injective. To do so, choose any $a_1, a_2 \in A$ where $a_1 \neq a_2$. We need to show that $f(a_1) \neq f(a_2)$.

We'll proceed by contradiction. Suppose that $f(a_1) = f(a_2)$. This means $f(f(a_1)) = f(f(a_2))$, which in turn tells us $a_1 = a_2$ because f is an involution. But that's impossible, since $a_1 \neq a_2$.

We've reached a contradiction, so our assumption was wrong. Therefore, we see that $f(a_1) \neq f(a_2)$, as required.

Proof Outline

1. Assume f is an involution.
2. Pick arbitrary $a_1, a_2 \in A$ such that $a_1 \neq a_2$.
3. Prove $f(a_1) \neq f(a_2)$.

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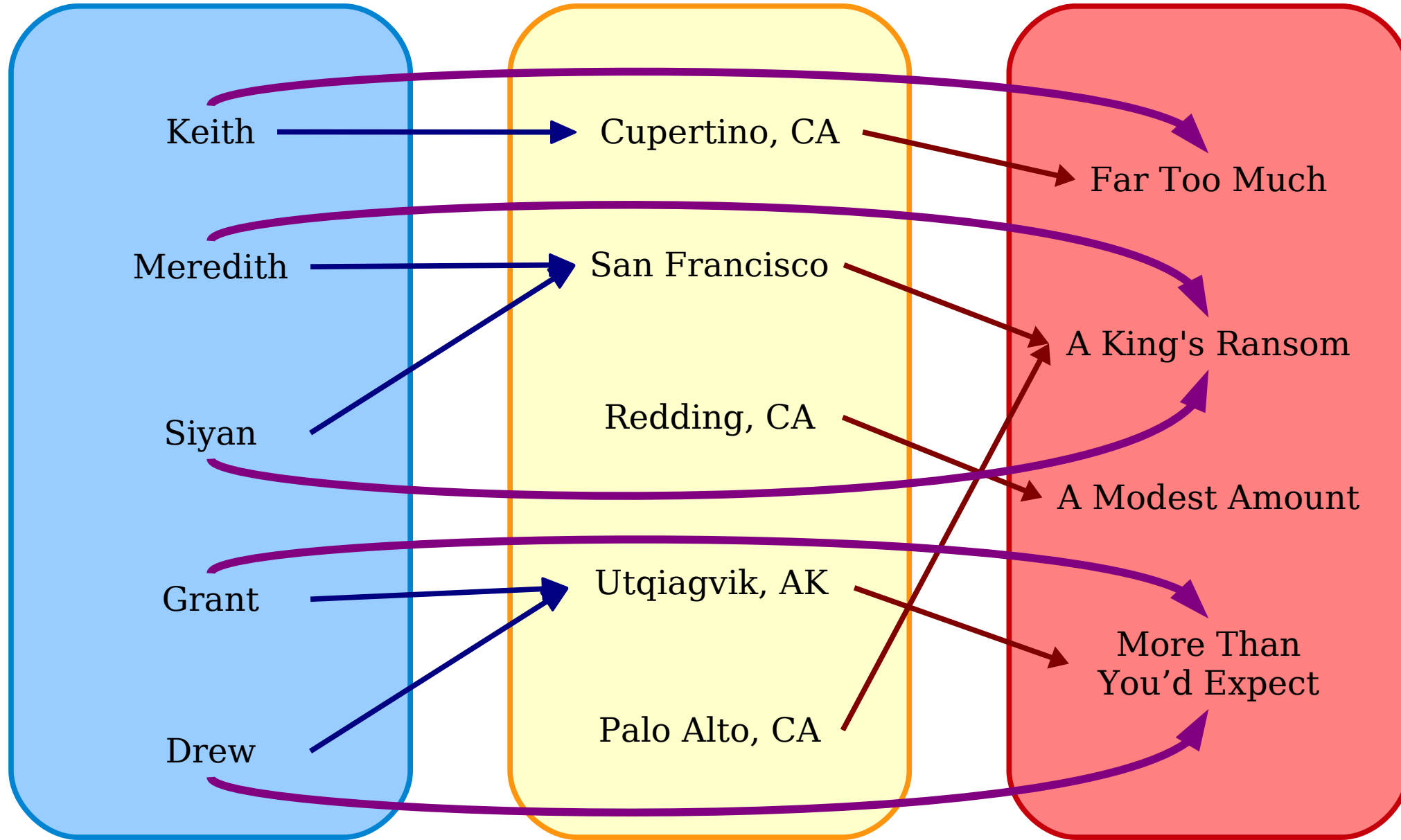
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1. Assume f is an involution.
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Function Composition

f : People → Places

g : Places → Prices



People

Places

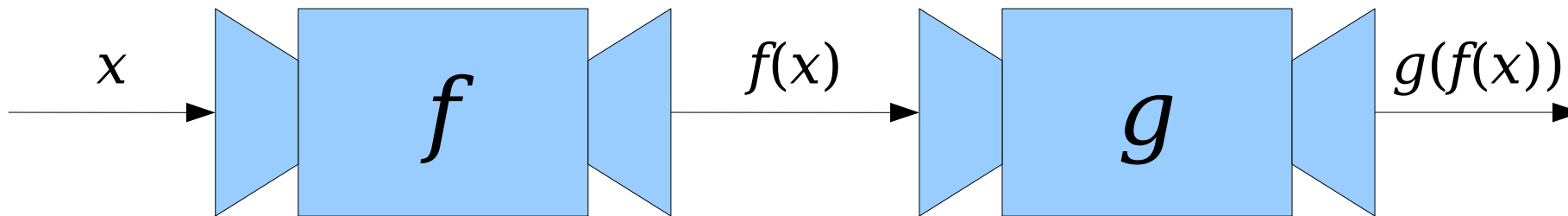
Prices

h : People → Prices

h(x) = g(f(x))

Function Composition

- Suppose that we have two functions $f : A \rightarrow B$ and $g : B \rightarrow C$.
- Notice that the codomain of f is the domain of g . This means that we can use outputs from f as inputs to g .



Function Composition

- Suppose that we have two functions $f : A \rightarrow B$ and $g : B \rightarrow C$.
- The **composition of f and g** , denoted **$g \circ f$** , is a function where
 - $g \circ f : A \rightarrow C$, and
 - $(g \circ f)(x) = g(f(x))$.
- A few things to notice:
 - The domain of $g \circ f$ is the domain of f . Its codomain is the codomain of g .
 - Even though the composition is written $g \circ f$, when evaluating $(g \circ f)(x)$, the function f is evaluated first.

The name of the function is $g \circ f$.
When we apply it to an input x ,
we write $(g \circ f)(x)$. I don't know
why, but that's what we do.

Properties of Composition

Theorem: If $f : A \rightarrow B$ is an injection and $g : B \rightarrow C$ is an injection, then the function $g \circ f : A \rightarrow C$ is an injection.

Organizing Our Thoughts

Theorem: If $f : A \rightarrow B$ is an injection and $g : B \rightarrow C$ is an injection, then the function $g \circ f : A \rightarrow C$ is an injection.

What We're Assuming

$f : A \rightarrow B$ is an injection.

$\forall x \in A. \forall y \in A. (x \neq y \rightarrow$
 $f(x) \neq f(y))$

$g : B \rightarrow C$ is an injection.

$\forall x \in B. \forall y \in B. (x \neq y \rightarrow$
 $g(x) \neq g(y))$

We're *assuming* these universally-quantified statements, so we won't introduce any variables for what's here.

What We Need to Prove

$g \circ f$ is an injection.

$\forall a_1 \in A. \forall a_2 \in A. (a_1 \neq a_2 \rightarrow$
 $(g \circ f)(a_1) \neq (g \circ f)(a_2))$

We need to *prove* this universally-quantified statement. so let's introduce arbitrarily-chosen values.

Theorem: If $f : A \rightarrow B$ is an injection and $g : B \rightarrow C$ is an injection, then the function $g \circ f : A \rightarrow C$ is an injection.

What We're Assuming

$f : A \rightarrow B$ is an injection.

$$\forall x \in A. \forall y \in A. (x \neq y \rightarrow f(x) \neq f(y))$$

$g : B \rightarrow C$ is an injection.

$$\forall x \in B. \forall y \in B. (x \neq y \rightarrow g(x) \neq g(y))$$

$a_1 \in A$ is arbitrarily-chosen.

$a_2 \in A$ is arbitrarily-chosen.

What We Need to Prove

$g \circ f$ is an injection.

$$\forall a_1 \in A. \forall a_2 \in A. (a_1 \neq a_2 \rightarrow (g \circ f)(a_1) \neq (g \circ f)(a_2))$$

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Now we're looking at
an implication. Let's
assume the antecedent
and *prove* the consequent.

Theorem: If $f : A \rightarrow B$ is an injection and $g : B \rightarrow C$ is an injection, then the function $g \circ f : A \rightarrow C$ is an injection.

What We're Assuming

$f : A \rightarrow B$ is an injection.

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$g \circ f$ is an injection.

$$\forall a_1 \in A. \forall a_2 \in A. (a_1 \neq a_2 \rightarrow (g \circ f)(a_1) \neq (g \circ f)(a_2))$$

Let's write this out separately and simplify things a bit.

Theorem: If $f : A \rightarrow B$ is an injection and $g : B \rightarrow C$ is an injection, then the function $g \circ f : A \rightarrow C$ is an injection.

What We're Assuming

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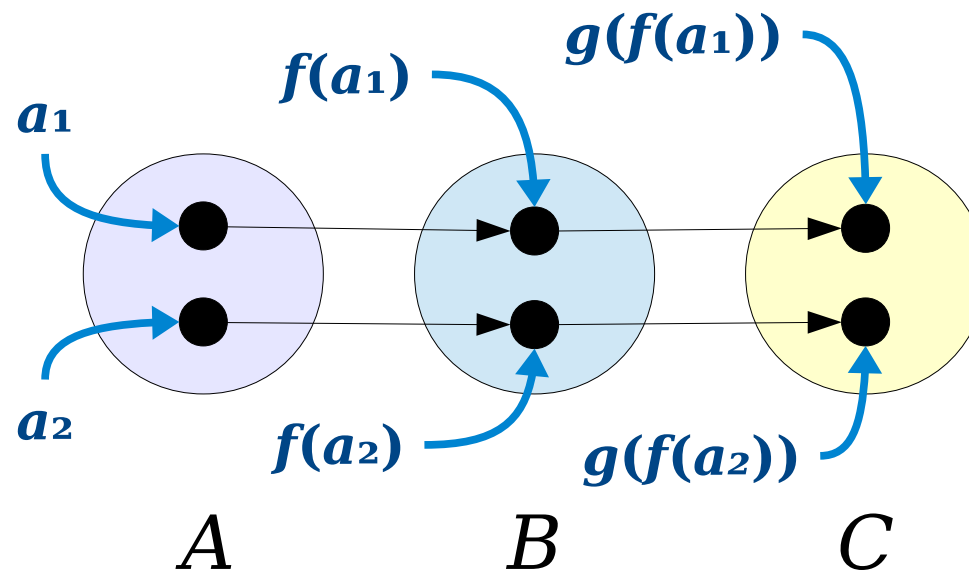
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What We Need to Prove

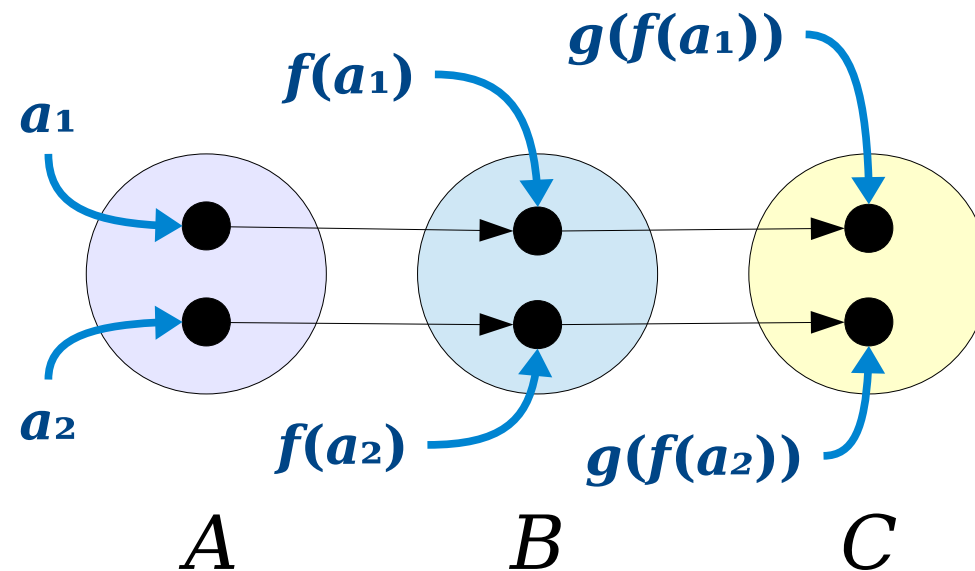
$g \circ f$ is an injection.

$$\forall a_1 \in A. \forall a_2 \in A. (a_1 \neq a_2 \rightarrow (g \circ f)(a_1) \neq (g \circ f)(a_2))$$

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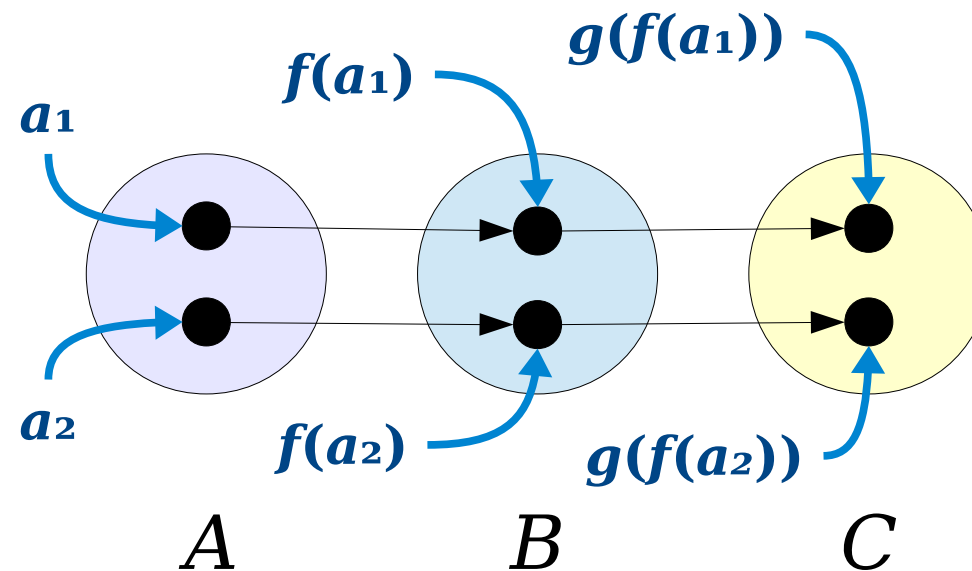


Theorem: If $f : A \rightarrow B$ is an injection and $g : B \rightarrow C$ is an injection, then the function $g \circ f : A \rightarrow C$ is also an injection.



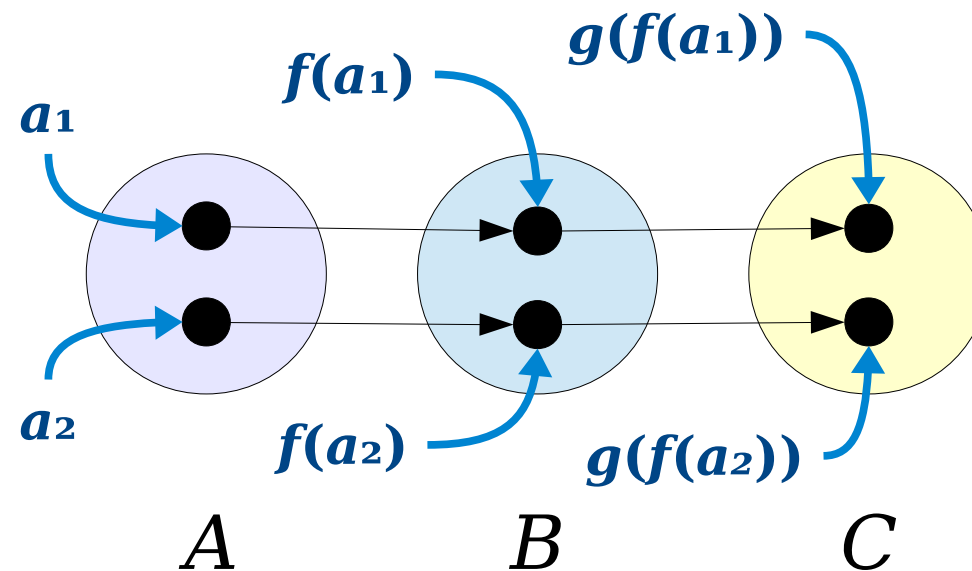
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Proof:



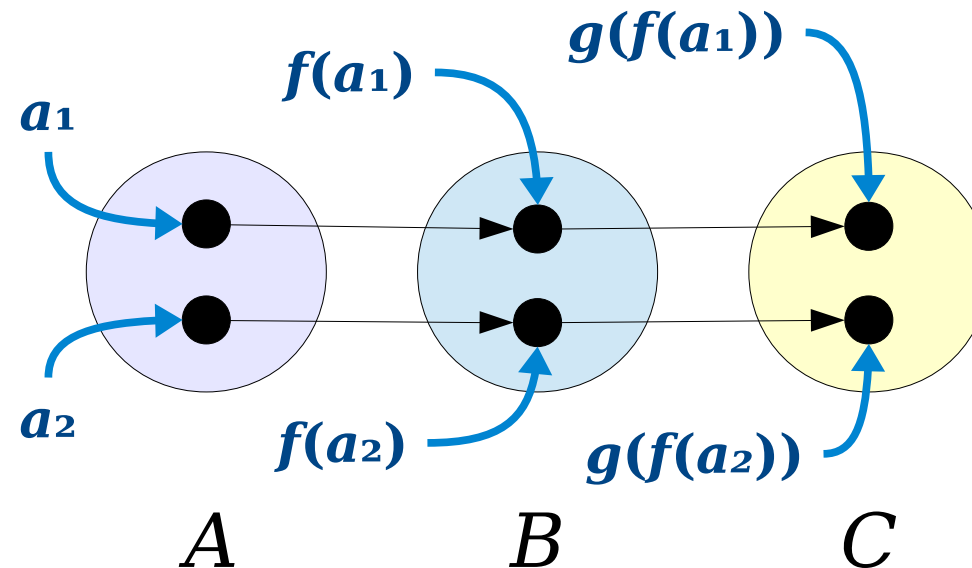
Theorem: If $f : A \rightarrow B$ is an injection and $g : B \rightarrow C$ is an injection, then the function $g \circ f : A \rightarrow C$ is also an injection.

Proof: Let $f : A \rightarrow B$ and $g : B \rightarrow C$ be arbitrary injections.



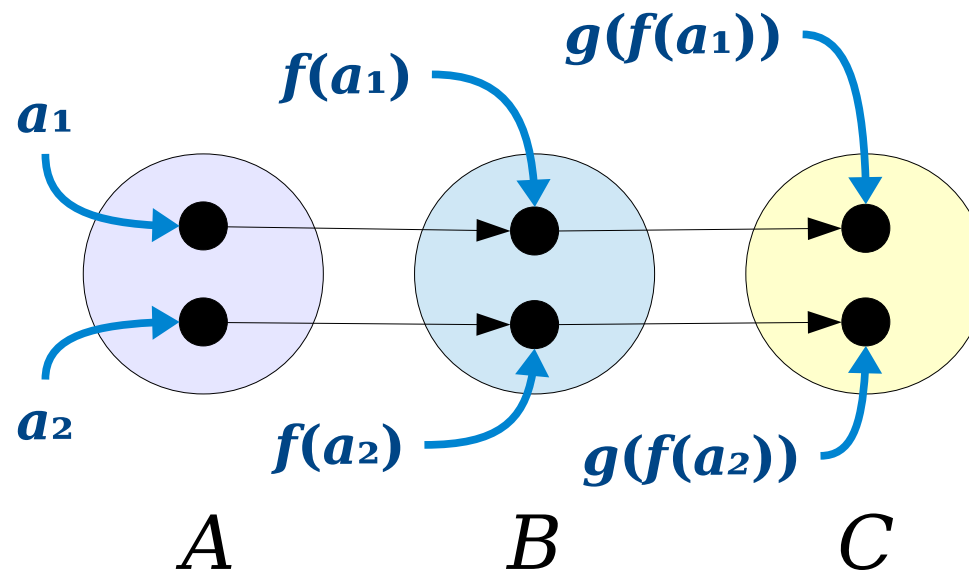
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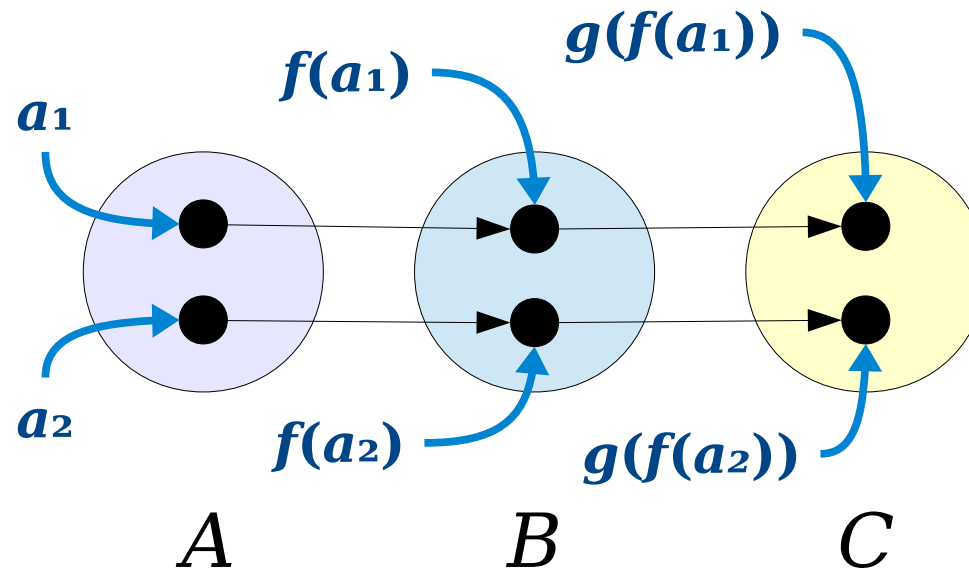
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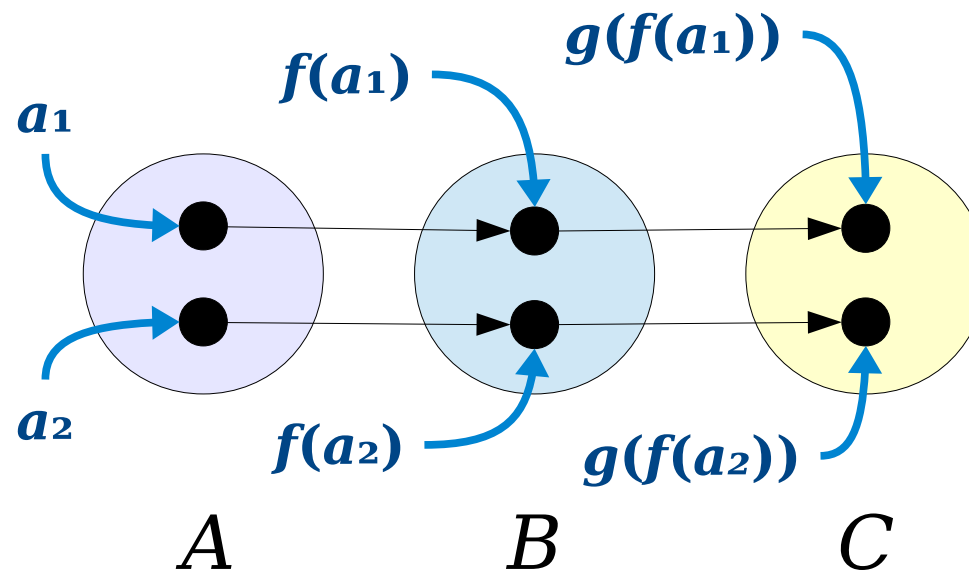
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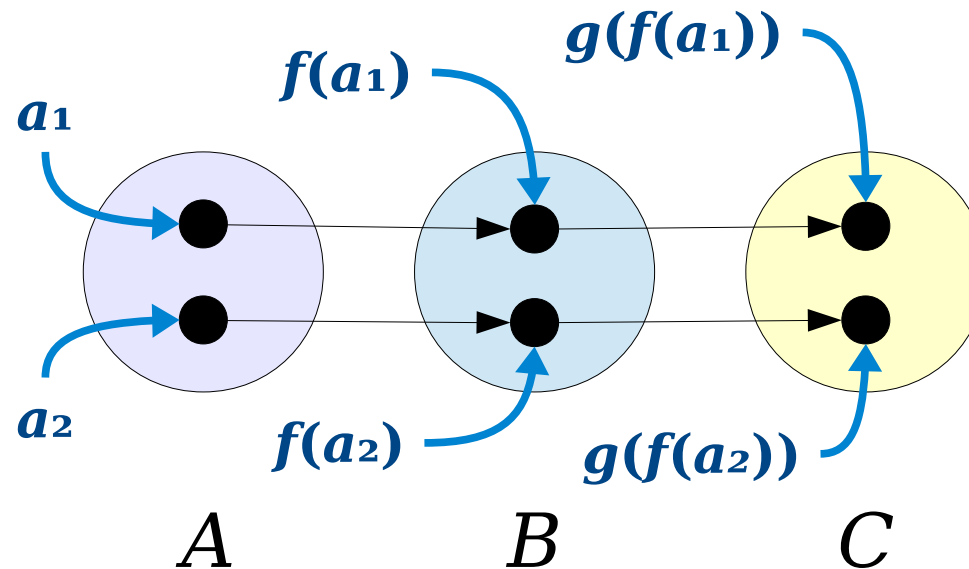
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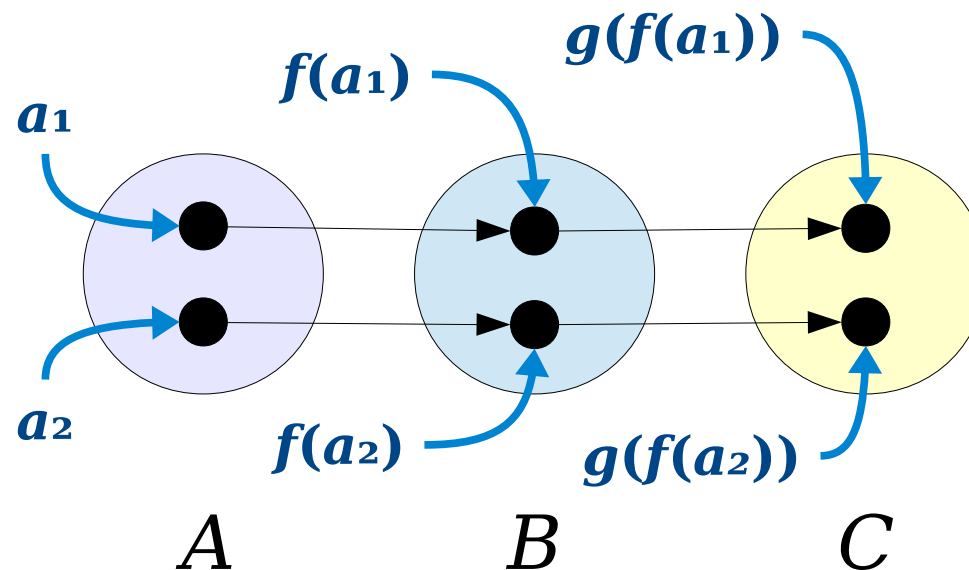
Since f is injective and $a_1 \neq a_2$, we see that $f(a_1) \neq f(a_2)$.



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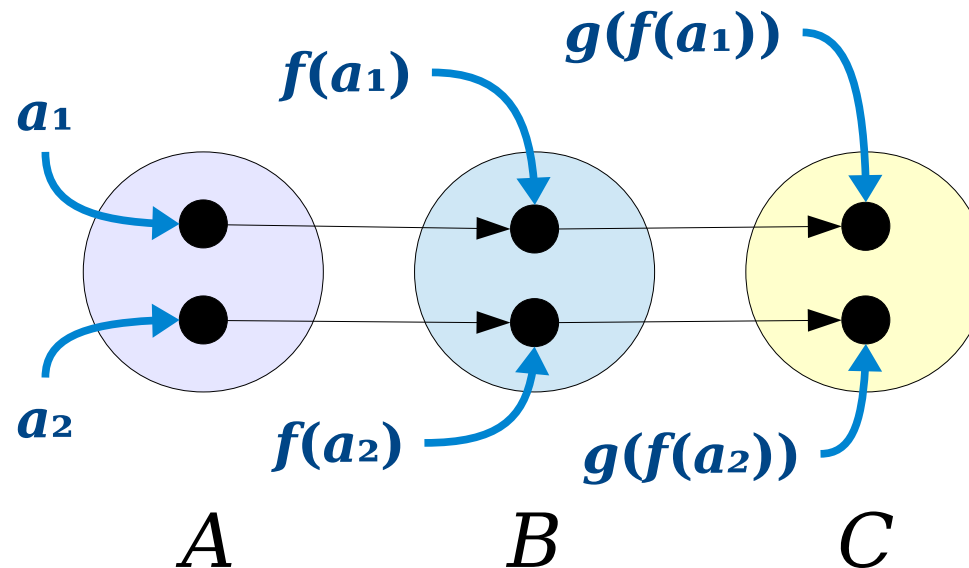
Since f is injective and $a_1 \neq a_2$, we see that $f(a_1) \neq f(a_2)$. Then, since g is injective and $f(a_1) \neq f(a_2)$, we see that $g(f(a_1)) \neq g(f(a_2))$, as required.



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Since f is injective and $a_1 \neq a_2$, we see that $f(a_1) \neq f(a_2)$. Then, since g is injective and $f(a_1) \neq f(a_2)$, we see that $g(f(a_1)) \neq g(f(a_2))$, as required. ■

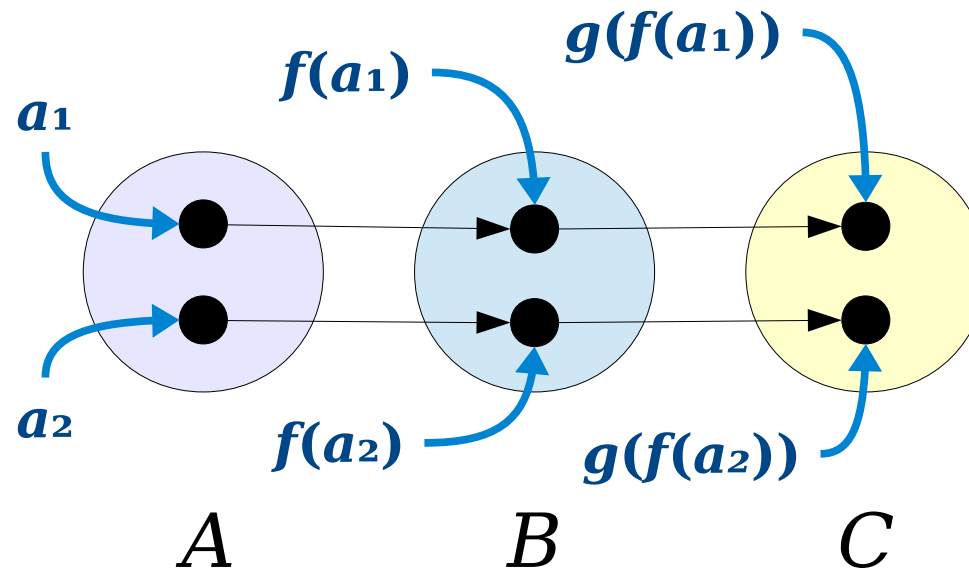


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Since f is injective and $a_1 \neq a_2$, we see that $f(a_1) \neq f(a_2)$. Then, since g is injective and $f(a_1) \neq f(a_2)$, we see that $g(f(a_1)) \neq g(f(a_2))$, as required. ■

Great exercise: Repeat this proof using the other definition of injectivity.



Major Ideas From Today

- Statements behave differently based on whether you're **assuming** or **proving** them.
- When you **assume** a universally-quantified statement, initially, do nothing. Instead, keep an eye out for a place to apply the statement more specifically.
- When you **prove** a universally-quantified statement, pick an arbitrary value and try to prove it has the needed property.
- As always: try concrete examples, draw pictures, etc. before you dive into writing a proof.

	To <i>prove</i> that this is true...	If you <i>assume</i> this is true...
$\forall x. A$	Have the reader pick an arbitrary x . We then prove A is true for that choice of x .	Initially, <i>do nothing</i> . Once you find a z through other means, you can state it has property A .
$\exists x. A$	Find an x where A is true. Then prove that A is true for that specific choice of x .	Introduce a variable x into your proof that has property A .
$A \rightarrow B$	Assume A is true, then prove B is true.	Initially, <i>do nothing</i> . Once you know A is true, you can conclude B is also true.
$A \wedge B$	Prove A . Then prove B .	Assume A . Then assume B .
$A \vee B$	Either prove $\neg A \rightarrow B$ or prove $\neg B \rightarrow A$. <i>(Why does this work?)</i>	Consider two cases. Case 1: A is true. Case 2: B is true.
$A \leftrightarrow B$	Prove $A \rightarrow B$ and $B \rightarrow A$.	Assume $A \rightarrow B$ and $B \rightarrow A$.
$\neg A$	Simplify the negation, then consult this table on the result.	Simplify the negation, then consult this table on the result.

Next Time

- ***Cardinality Revisited***
 - Formalizing our definitions.
- ***The Nature of Infinity***
 - Infinity is more interesting than it looks!
- ***Cantor's Theorem Revisited***
 - Formally proving a major result.

Extra Slides

(The following is a proof of a theorem just like the one we just did with injection, but with surjection.)

Theorem: If $f : A \rightarrow B$ is a surjection and $g : B \rightarrow C$ is a surjection, then the function $g \circ f : A \rightarrow C$ is a surjection.

Theorem: If $f : A \rightarrow B$ is surjective and $g : B \rightarrow C$ is surjective, then $g \circ f : A \rightarrow C$ is also surjective.

Theorem: If $f : A \rightarrow B$ is surjective and $g : B \rightarrow C$ is surjective, then $g \circ f : A \rightarrow C$ is also surjective.

Proof:

Theorem: If $f : A \rightarrow B$ is surjective and $g : B \rightarrow C$ is surjective, then $g \circ f : A \rightarrow C$ is also surjective.

Proof: Let $f : A \rightarrow B$ and $g : B \rightarrow C$ be arbitrary surjections.

Theorem: If $f : A \rightarrow B$ is surjective and $g : B \rightarrow C$ is surjective, then $g \circ f : A \rightarrow C$ is also surjective.

Proof: Let $f : A \rightarrow B$ and $g : B \rightarrow C$ be arbitrary surjections. We will prove that the function $g \circ f : A \rightarrow C$ is also surjective.

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What does it mean for $g \circ f : A \rightarrow C$ to be surjective?

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$$\forall c \in C. \exists a \in A. (g \circ f)(a) = c$$

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Therefore, we'll choose an arbitrary $c \in C$ and prove that there is some $a \in A$ such that $(g \circ f)(a) = c$.

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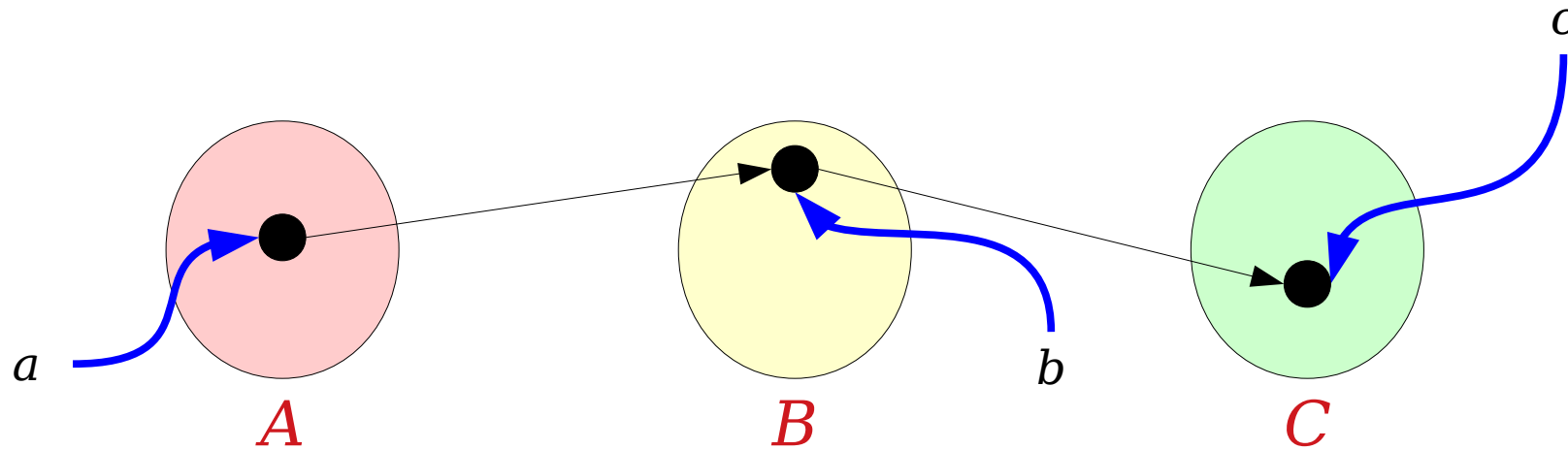
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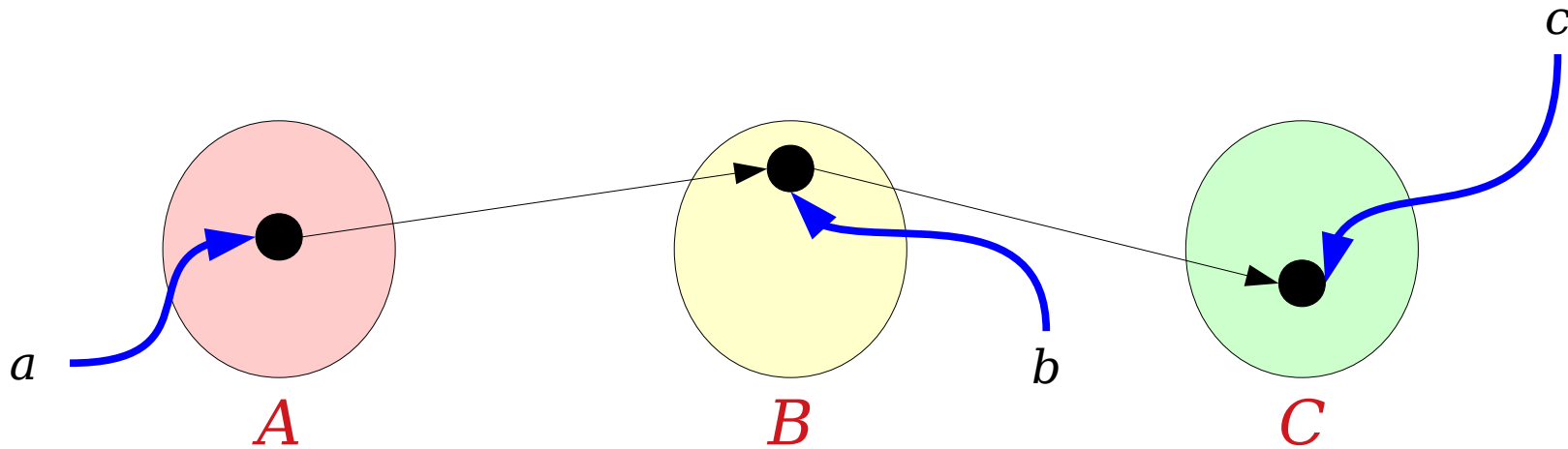
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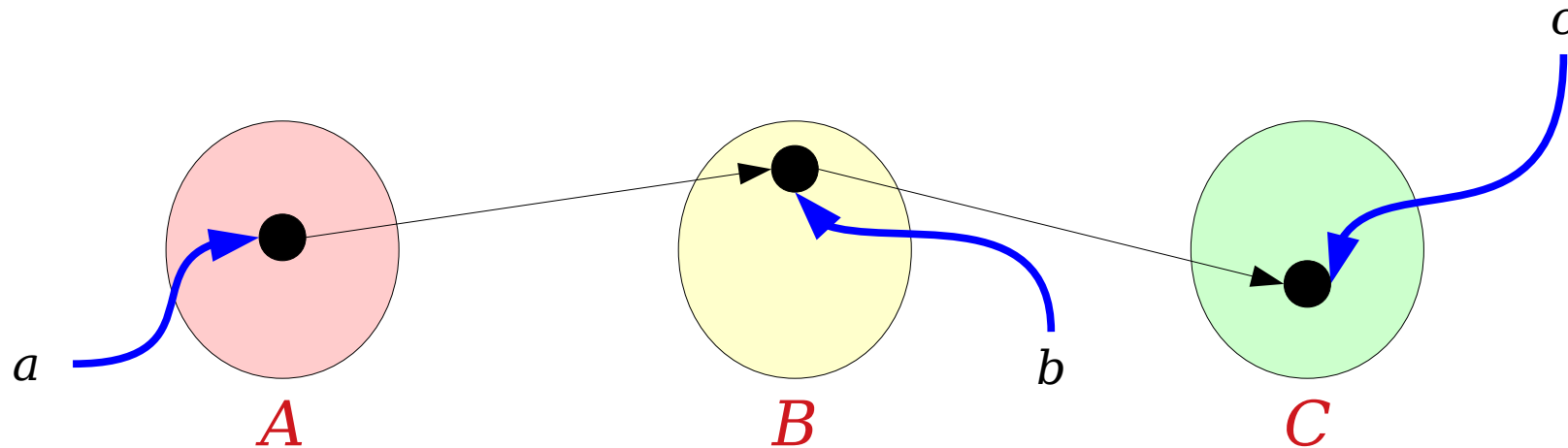
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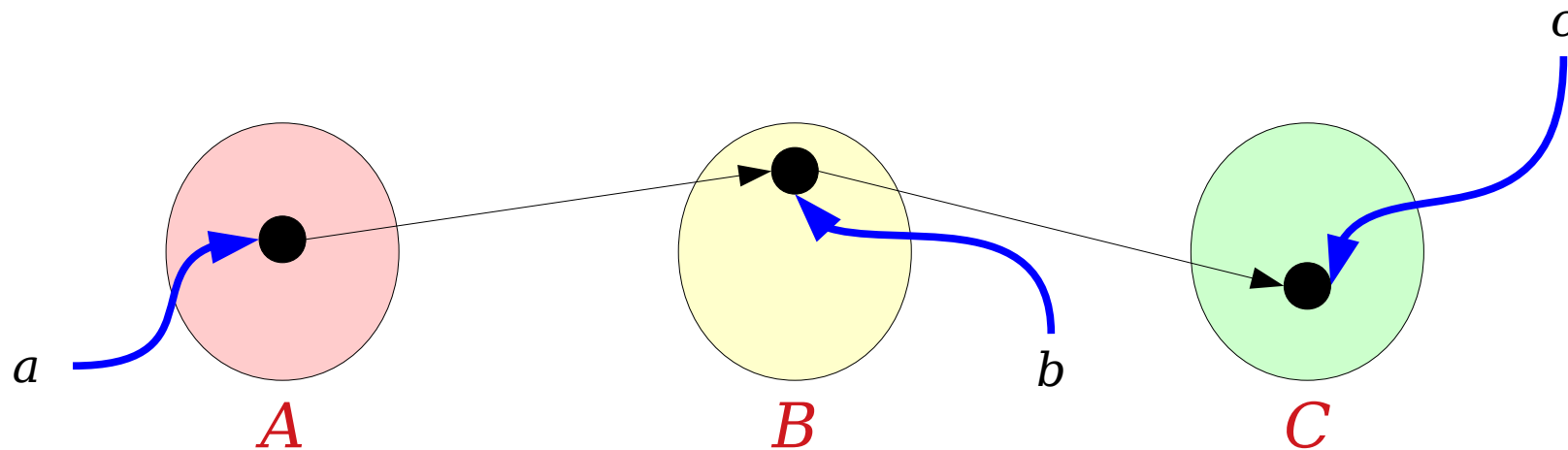
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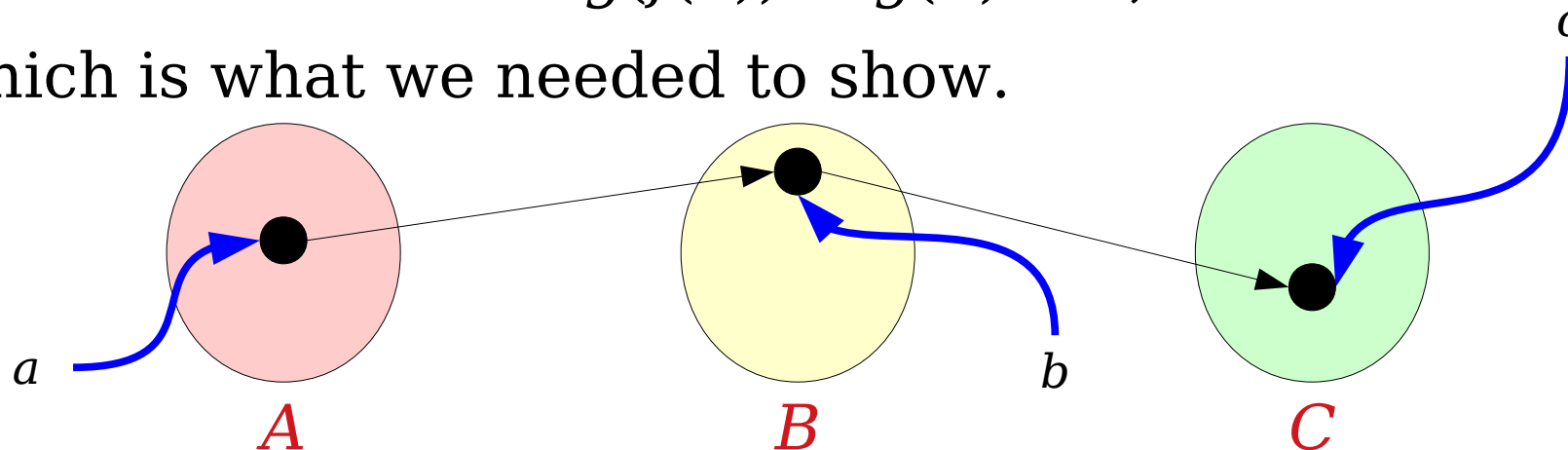
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